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중요표본추출 기법 이용한  
重要標本抽出 技法 利用한

交通網 區間의 混雜確率 推定  
교통망 구간 부수 확률 관측

ESTIMATION OF THE CONGESTION PROBABILITY ON A TREE-TYPE  
TRANSPORTATION NETWORK BY IMPORTANCE SAMPLING

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要 約

본 논문의 목적은 교통망 분석에 있어서 중요한 그러나 혼회 발생하지 않는 사건의 발생확률을 추정하는 방법론 개발에 있다. 예를 들어, 안정적(stable) 교통망에서 일시적인 혼잡현상이 발생할 확률을 씨플레이손을 통해 추정하는 방법에 관한 것이다. 이 분야에 활발한 연구([3], [12])가 있어 왔으나 개괄적(Heuristic)방법에 제한되어 있었다.

본 논문은 위 문제에 대하여 포괄적(unified)이며 이론적인 방법론을 제시하였다. 이를 위해 대 분산이론(Large Deviation Theory)과 중요표본추출(Importance Sampling)기법이 이용되었으며 예로서 사용된 망은 두개의 구간이 이어진 교통망이다. 부수적으로 혼잡현상의 가장 대표적 형태를 구하는 방법이 제시되었다.

I . Introduction

Consider a stable system of a tree-type network of queues. Let  $\{X_n\}$  denote the embedded r-dimensional Markov chain representing the system. Our goal is to estimate

the probability that  $\text{Max}_{1 \leq i \leq r} X_{ni}$  or  $\sum_{1 \leq i \leq r} X_{ni}$  reaches a large number K within a busy cycle by simulation.

Direct simulation is very expensive, inefficient, and even impossible in cases, since the events of interest are extremely rare.

Instead, the change of measure for importance sampling is considered. It is known that if a Markov chain obeys a large deviation principle(LDP) with a rate function, the efficient change of measure for importance sampling can be obtained successfully.

After reviewing the use of importance sampling in simulating a M/M/1 queue, the paper apply this idea to the case of two queues in tandem by providing the LDP of the embedded Markov chain of the system.

II . Review of M/M/1 Queue and Importance Sampling

Before stating the new results, a brief summary of the relevant theory relating to M/M/1 queue will be given. Consider a stable M/M/1 system,  $\bar{S}$ , with arrival rate  $\mu$  and service rate  $\lambda$  with  $\mu/\lambda < 1$ . We are interested in estimating  $\alpha = \Pr[A]$ , where  $\{X_n\}$  is a sequence of successive population in the queue and A denotes the event that  $X_n$  hits K before becoming 0, assuming it is initially zero. Our concern is to understand how these values can be obtained by efficient simulation.

The estimated value of  $\alpha$ ,  $\hat{\alpha}$ , may be obtained by direct simulation as follows:

Let us call a "cycle" a movement from 0 to the first time either 0 is reached again, or K is reached. Define  $V_m = 1$  if  $X_n$  reaches K in the  $m^{th}$  cycle. Since  $\{X_n\}$  is a regenerative process with a regenerating point=0,  $V_m$ 's are independent and identically distributed with  $E[V_m] = \alpha$ . Therefore,  $\hat{\alpha} = \frac{V_1 + \dots + V_p}{p}$  is an unbiased estimator of  $\alpha$ .

To guarantee the following accuracy.

$$Pr\left[ \frac{|\hat{\alpha} - \alpha|}{\alpha} < a \right] > b$$

we have to generate  $N\bar{s}$  cycles

$$N_s) \frac{\gamma(b)^2 \text{Var}(V_m)}{a^2 \alpha^2}$$

where  $P[N(0,1) < \gamma(b)] = b$ .

Since we are interested in a very small  $\alpha$ ,  $\text{Var}(V_m) = \alpha - \alpha^2 \sim \alpha$ . Therefore,

$$N_s) \frac{\gamma(b)^2}{a^2 \alpha}$$

It shows that the minimum  $N\bar{s}$  depends on  $\alpha$ . Since  $\alpha$  has a very small value, it implies that we have to generate a large number of cycles to obtain precision to some degree. It is the main source of inefficiency in direct simulation.

To overcome this inefficiency, we consider "Importance Sampling". Main idea of importance sampling is to modify the original system so that its population grows much faster and, in consequence, the event happens more frequently and translate the estimate from the modified system in terms of the original one. For example, if a modification S of the original  $\bar{S}$  is defined as M/M/1 with arrival rate  $\mu$  and service rate  $\lambda$ , the event is more frequent in S than in  $\bar{S}$ .

Suppose  $P\bar{s}$ ,  $P_s$  are the probability measures that govern  $\bar{S}$ , S, respectively and  $L_m$ 's are the likelihood ratio between  $P_s$  and  $P\bar{s}$  in the  $m^{th}$  cycle. ( $L_m$ 's can be easily computed for the  $m^{th}$  cycle.) Since the event is more likely under  $P_s$  than under  $P\bar{s}$ ,  $L_m > 1$  whenever  $V_m = 1$ . Therefore,  $\text{Var}_s(L_m V_m) < \text{Var}_{\bar{s}}(V_m)$ . It implies that, to obtain a same error bound, we have to generate more cycles for  $\bar{S}$  than for S.

Therefore, the center of importance sampling is on the question "what is the modifi-

cation that minimizes  $L_m$  when  $V_m=1$ ?" In other words, we have to find "the most likely way" that  $S$  reaches  $K$  and then modify the system to make the most likely way exceedingly likely. In case of  $M/M/1$ , there are various methods to find the most likely way. (cf. [3], [12]).

Unfortunately, it does not seem to be possible to apply those methods used here to more complicated examples, such as the system of a pair of queues in tandem and a 3 dimensional tree-type network, etc.

### III . Markov Chain and Large Deviation Theory

In this section, we present a large deviation theorem due to Wentzel[8], regarding certain Markov chains. Consider the parameterized Markov chain  $\{X_n^\epsilon\}$  given by

$$X_0^\epsilon = x_0$$

$$X_{n+1}^\epsilon = X_n^\epsilon + \epsilon V(X_n^\epsilon, \zeta_n),$$

where  $\epsilon > 0$  is the parameter which defines the Markov chain  $\{X_n^\epsilon\}$ ,  $x_0$  is the initial value,  $V(\cdot, \cdot)$  is a function from  $R^1 \times R^1$  to  $R^1$  and  $\zeta_n$ 's are i. i. d. r. v. 's. We are interested in analyzing  $\{X_n^\epsilon\}$  when  $\epsilon \rightarrow 0$ .

Let  $F_x$  denote the d. f. of  $V(x, \zeta_n)$ .

$$\text{Let } m(x) = \int_{R^1} z dF_x(z)$$

$$M_x(s) = \int_{R^1} \exp\langle s, z \rangle dF_x(z)$$

$$l_x(s) = \log M_x(s)$$

$$h_x(u) = \sup_{s \in R^1} [\langle s, u \rangle - l_x(s)].$$

Assume the following :

- (a)  $M_x(s) < \infty$  in a neighborhood of 0 for each  $x \in R^1$
- (b)  $d(F_{x_1}, F_{x_2}) < c | X_1 - X_2 |$ , where  $d$  is the Prohorov distance (see [1]),  $c > 0$  is a constant and  $| \cdot |$  is a Euclidean norm.

Next, construct continuous-time paths from the realization of  $\{X_n^\epsilon\}$ . To do this, at the epochs  $t = n \cdot \epsilon$ , let  $X^\epsilon(t) = X_n^\epsilon$  and interpolate piecewise linearly. Let  $C_T$  denote the set of the piecewise continuously differentiable functions  $\phi : [0, T] \rightarrow R^1$  such that  $\phi(0) = x_0$  is fixed. Let  $P^\epsilon$  denote the measure induced by the  $\{X^\epsilon(t)\}$  on the Borel-field of  $C_T$ , endowed with the Skorohod topology [1].

Define the rate function

$$I(x_0, \phi) = \begin{cases} \int_0^T h_{\phi(t)}(\phi'(t)) dt, & \phi(0) = x_0 \\ \infty & , \text{ o. w. . .} \end{cases}$$

Under (a), (b) with a couple of more technical assumptions we have the following results.

Theorem 1.

i ) For each closed subset  $F$  of  $C_T$ ,

$$\limsup_{\epsilon \rightarrow 0} -\epsilon \log P_x^\epsilon(F) \leq \inf_{\Psi \in F} I(x, \Psi).$$

ii ) for each open subset  $G$  of  $C_T$ ,

$$\liminf_{\epsilon \rightarrow 0} -\epsilon \log P^\epsilon(G) \geq \inf_{\Psi \in G} I(x, \Psi).$$

Therefore, if  $\inf_{\Psi \in A} I(\Psi) = \inf_{\Psi \in A^c} I(\Psi)$ ,

then  $P_x^\epsilon(A) \sim \exp[-\frac{1}{\epsilon} \inf_{\Psi \in A} I(x, \Psi)]$  (UTLE), as  $\epsilon \rightarrow 0$ .

(where UTLE is the acronym for up to logarithmic equivalence.)

Suppose that we want to estimate the probability of event  $A$ ,

$A = \{\text{Max}_t X_n$  exceeds  $K$  before hitting 0, given  $X_0 = x_0\}$ .

We consider the importance sampling method for this purpose.

Our goal is to find the optimal modification  $S^*$ , which satisfies

$\text{Var } s^*(Ls^*mV_m) < \text{Var } s(Ls mV_m)$ , for all modification  $S$ .

In other words, we want to solve the fol-

lowing optimization problem.

$$\text{Min}_S \int_A L_s^2 dP_s,$$

when  $L_s = \frac{dP_s}{dP_s}$  is the likelihood ratio.

The following two observations are critical to find the optimal  $S^*$ ;

Observation 1

If  $S^*$  is the optimal solution of the problem, then it also minimize

$$\int_A L_s^2 dP_s^*$$

Observation 2

If  $\phi^*$  satisfies

$$\inf_{\phi \in C_T} I_s(x, \phi) = I_s(x, \phi^*)$$

then  $S^*$  solves the following equation

$$I_s^*(x, \phi^*) = 0$$

Then, for any arbitrary small  $\delta > 0$ ,

$P_s^*(\phi_s^*)$  is asymptotically equal to 1.

as  $\epsilon \rightarrow 0$ , where  $\phi_s^* = \{\phi \in C_T : \|\phi - \phi^*\| < \delta\}$ . ( $\|\cdot\|$  is a uniform norm on  $C_T$ .)

In words,  $\phi^*$  is the most likely path of the event A and  $S^*$  is the system in which  $\phi^*$  is the most likely path among all the sample paths.

When we consider the Markov chain of queueing system  $\{X_n\}$ , the distribution function of the probability measure P, governing  $\{X_n\}$ , violates the assumed condition (b). Hence, we can not blindly use these two observation to the case of queueing process in finding the optimal modification for importance sampling.

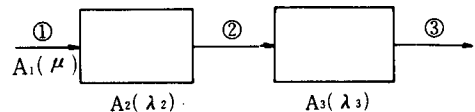
The discontinuity in the probability structure of a queueing system (violation of condition (b)), is mainly caused by the nonnegativity condition on the queue length. If there is no customer in a station, the station is in an idle state and no service is offered. This can be thought of a nonnegativity constraint on the queue length pro-

cess, which imposes a change in the probabilistic behavior of the server when its queue becomes empty. This is the main obstacle in finding the LDP of the queueing system through the above results.

#### IV. Potential Process and its LDP

In order to avoid the nonnegativity constraint, we present potential process and Skorohod problem regarding networks of tandem queues, due to Park [9].

Let's Consider a network of tandem queues.



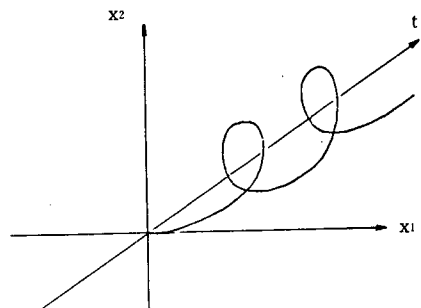
(Figure 1) Network of Tandem Queues (( $\mu, \lambda_1, \lambda_2$ )-System)

$A_i(t)$ : Arrival Process ~ Poisson Process with  $\mu$ .

$A_i(t)$ : Service time ~ Exponential dist. with  $\lambda_i$  and  $\mu < \lambda_i$ , for  $i=1, 2$  (stable system).

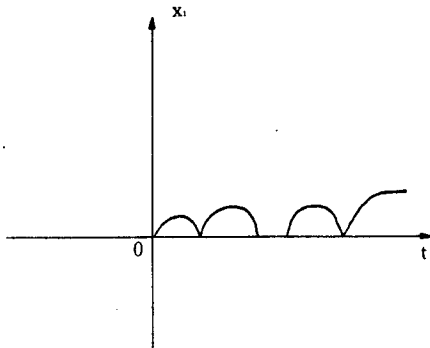
$$X(t) = [X_1(t) X_2(t)],$$

where  $X_i(t)$  is the number of customers at the station  $i$  at time  $t$ .  $X(t)$  is called the queue length process of tandem queues.

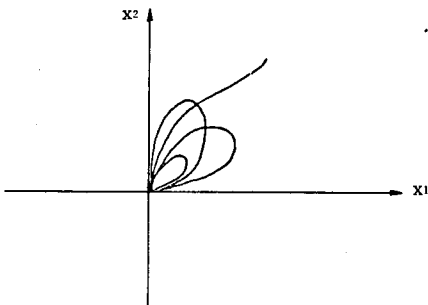


(Figure 2) A Sample Trajectory on  $(t, x_1, x_2)$

- Busy cycle is a movement from the empty state to the next empty state.
- Each cycle on the figure 4 represents a busy cycle. (Empty state part of trajectory is hidden at the origin.)



〈Figure 3〉 A Sample Trajectory in Station i.



〈Figure 4〉 A Sample Trajectory on  $(x_1, x_2)$

We consider the potential process described as follows;  
 Let's consider a system in which servers provide units of services and they are not concerned whether there are customers in their stations or not. The time needed for the server in the station  $i$  generating a unit of service follows the exponential distribution with parameter  $\lambda_i$ . If they provide a service unit at a time when their queues are empty, then we consider the

“queue length” to change from 0 to  $-1$ . Similarly if the queue length is already negative we simply decrement it by 1 again at the service time. The rest of the topological structure and probabilistic assumptions are defined same as the one for the original system. The variable we observe at time  $t$  at each station is the number of arrivals up to time  $t$  to the station minus the number of the service units generated in the station up to time  $t$ . Let  $Y_i(t)$  be the observation in  $i$ th element of  $Y(t)$ . We call  $Y$  “potential process”.

Based on the results in the section 3, LDP of the Markov chain of the potential process,  $\{Y_n\}$ , can be easily established and its rate function on  $[0, T]$  is

$$I(x, \phi) = \begin{cases} \int_0^t h_x(\phi'(t)) dt, & \text{if } \phi(0) = x \\ \infty & , \text{o. w.} \end{cases}$$

$$\text{where } h_x(U) = \inf_{\substack{U_1 = S_1 - S_2 \\ U_2 = S_2 - S_3}} \left\{ s_1 k_1 \left( \frac{1}{s_1} \right) + s_2 k_2 \left( \frac{1}{s_2} \right) + s_3 k_3 \left( \frac{1}{s_3} \right) \right\}$$

$$\text{and } k_v(s) = \begin{cases} ys - \log ys - 1, & s > 0 \\ \infty & , \text{o. w.} \end{cases}$$

When we calculate  $I(x, \phi)$ , we assume  $\mu + \lambda_1 + \lambda_2 = 1$ , without loss of generality.

It is noticed that  $h_x(\cdot)$  does not depend on  $x$ . Therefore, we denote  $h_x(\cdot)$  with  $h(\cdot)$ .

Therefore, LDP of  $\{Y_n\}$  implies that

$$P_{\epsilon}^t(A) \sim \text{Exp} \left[ -\frac{1}{\epsilon} \inf_{\phi \in A} \int_0^t h(\phi'(t)) dt \right],$$

$\epsilon \rightarrow 0$ , where  $T_{\epsilon}$  is the first time when  $\text{Max } \phi_i$  reaches  $K$ .

V . Skorohod Problem

Our objective in this section is to exhibit the queue length process of a network of tandem queues as a continuous mapping Z of the potential process Y, i.e.,  $Z = \theta \cdot Y$ . Then Z should have the same distribution as the queue length process of the network of the tandem queues.

The following theorem is a key to achieve the goal.

Let  $D^2_T$  be a set of 2 dimensional right continuous functions with left limits for all t on [0, T].

$$\text{Let } M = \begin{bmatrix} 1 & -1 \\ 0 & 1 \end{bmatrix}$$

Theorem 2.

For a given  $\Psi(\cdot)$  with  $\Psi(\cdot) \in R^2(t)$ , there is a unique pair of function  $(\phi, u)$ , satisfying the following :

- (1)  $\phi(t) = \Psi(t) - Mu(t)$
- (2)  $\phi(t) \geq 0$ .
- (3)  $u_i(\cdot)$  is non-decreasing with  $u_i(0) = 0$  and  $u_i(\cdot)$  increases only at those times t where  $\phi_i(t) = 0$

For a proof, see Park[9].

The problem of showing the existence of  $\phi, u$ , is called "Skorohod problem" and the pair  $(\phi, u)$  is called the solution of the Skorohod problem.

Let  $\theta$  and  $\sigma$  be a mapping from  $C_T$  to  $C_T$ , defined by  $\theta(\Psi) = \phi$  and  $\sigma(\Psi) = u$ .

We name  $\theta$  "Skorohod map". Based on this map.

Let  $Z(\omega)(t) = \theta(Y(\omega))(t)$ ,  $U(\omega) = \sigma(Y(\omega))(t)$ . In Lemma 2, we investigate the probability structure of the process Z.

Lemma 2

- (1) The processes Z and U satisfy the fol-

lowing almost surely;

- 1.  $Z(t) \geq 0 \in R^2$ .
  - 2.  $U_i(t)$  is nondecreasing with  $U_i(0) = 0$ .
  - 3.  $U_i(t)$  increases only at those time t where  $Z_i(t) = 0$ ,
- (2) Z is a time homogeneous Markov process.

It implies that the process Z and the queue length process X have the same small time increment conditional probability and, therefore, they have the same infinitesimal generator. It follows that they have the same distribution. (cf. p.111 and p. 161 in Ethier and Kurtz[7].)

VI . LDP of Network of Tandem queues

Since the process Z and the queue length process X have the same distribution, the Markov chain  $\{Z_n\}$  of Z and  $\{X_n\}$  have the same probability structure. It follows that for any closed set F,

$$P^{x_0} \{Z_n \in F\} = P^{x_0} \{X_n \in F\}$$

and  $\theta^{-1}(F)$  is a closed set.

Therefore, by the property of the LDP of the potential process,

i) for any closed set F,

$$\limsup_{\epsilon \rightarrow 0} -\epsilon \log P^{x_0} \{Z_n \in F \mid x_0 = x\} \leq \inf_{\phi \in F} \inf_{\Psi \in \theta^{-1}(\phi)} I(x, \Psi)$$

By the same reason.

ii) for any open set G,

$$\liminf_{\epsilon \rightarrow 0} -\epsilon \log P^{x_0} \{Z_n \in G \mid x_0 = x\} \geq \inf_{\phi \in G} \inf_{\Psi \in \theta^{-1}(\phi)} I(x, \Psi)$$

Therefore, LDP of  $\{X_n\}$  is well-defined with its rate function;

$$I_{\{X_n\}}(x, \phi) = \begin{cases} \inf_{\Psi \in \theta^{-1}(\phi)} \int_0^t h(\Psi(t)) dt, & \text{if } \phi(0) = x \\ \infty, & \text{o. w. .} \end{cases}$$

**VI. The Most Efficient Change of Measure for Importance Sampling**

The observation 1 and 2 present the method to find the optimal change of measure.

First, we find the most likely path of the event A.

$$\begin{aligned} & \inf_{\phi \in A} I|X_n| (x, \phi) \\ &= \inf_{\phi \in A} \inf_{\Psi \in \mathcal{G}^{-1}(\phi)} I|Y_n| (\Psi), \\ & \text{where } A_0 = \{ \phi \in A : \phi(0) = 0 \} \\ &= \inf_{\phi \in A} \inf_{\Psi \in \mathcal{G}^{-1}(\phi)} \int_0^{T_\phi} h(\Psi(t)) dt \\ &= \inf_{\substack{S_1 > 0, S_2, S_3 > 0 \\ S_1 - S_2 > 0 \text{ or } S_2 - S_3 > 0 \\ S_1 + S_2 + S_3 = 1}} T_\phi \left[ s_1 h_{\lambda_1} \left( \frac{1}{S_1} \right) + s_2 h_{\mu_1} \left( \frac{1}{S_2} \right) + s_3 h_{\mu_2} \left( \frac{1}{S_3} \right) \right], \end{aligned}$$

$$\text{where } T_\phi = \begin{cases} \frac{1}{S_1 - S_2}, & \text{if } S_1 < S_2 \text{ and } S_2 < S_3 \\ \frac{1}{S_1 - S_2} & \text{o. w.} \end{cases}$$

This nonlinear programming problem can be easily solved by Kuhn-Tucker condition and the minimum is achieved at  $S_1 = \lambda_1$ ,  $S_2 = \mu$ ,  $S_3 = \lambda_2$ , when  $\lambda_1 < \lambda_2$  and at  $S_1 = \lambda_1$ ,  $S_2 = \lambda_2$ ,  $S_3 = \mu$ , when  $\lambda_2 < \lambda_1$

That is, the minimum is obtain when  $\mu$  is exchanged with smaller of  $\lambda_1$  and  $\lambda_2$

**VII. Future Research**

This result can be easily generalized to the case of general tree-type networks. To study the congestion event in transportation networks more precisely, this research should be generalized to the case that the service rates and arrival rates are functions of state variables. The focus of continuing effort will be in this direction.

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