

## A NOTE ON FINITE CONDITIONS OF ORTHOMODULAR LATTICES

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ABSTRACT. We prove the following: every chain-finite OML is path-connected; every finite block of an OML  $L$  is path-connected with at least one other block in  $L$ ; every OML with uniformly finite sites is path-connected.

### 1. Preliminaries

Several finite conditions of orthomodular lattices have been investigated [3, 6]. In this paper, we prove some properties of orthomodular lattices with some finite conditions.

An *orthomodular lattice* (abbreviated by OML) is an ortholattice  $L$  which satisfies the *orthomodular law*: if  $x \leq y$ , then  $y = x \vee (x' \wedge y)$  [5]. A *Boolean algebra*  $B$  is an ortholattice satisfying the *distributive law*:  $x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z) \quad \forall x, y, z \in B$ .

A *subalgebra* of an OML  $L$  is a nonempty subset  $M$  of  $L$  which is closed under the operations  $\vee$ ,  $\wedge$  and  $'$ . We write  $M \leq L$  if  $M$  is a subalgebra of  $L$ . If  $M \leq L$  and  $a, b \in M$  with  $a \leq b$ , then the *relative interval sublattice*  $M[a, b] = \{x \in M \mid a \leq x \leq b\}$  is an OML with the *relative orthocomplementation*  $\#$  on  $M[a, b]$  given by  $c^\# = (a \vee c') \wedge b = a \vee (c' \wedge b) \quad \forall c \in M[a, b]$ . In particular,  $L[a, b]$  will be denoted by  $[a, b]$  if there is no ambiguity.

The *commutator* of  $a$  and  $b$  of an OML  $L$  is denoted by  $a * b$ , and is defined by  $a * b = (a \vee b) \wedge (a \vee b') \wedge (a' \vee b) \wedge (a' \vee b')$ . The set of all commutators of  $L$  is denoted by  $ComL$  and  $L$  is said to be *commutator-finite* if  $|ComL|$  is finite [2, 4]. For elements  $a, b$  of an OML, we say  $a$  *commutes with*  $b$ , in symbols  $a \mathbf{C} b$ , if  $a * b = 0$ . If  $M$  is a subset of an OML

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$L$ , the set  $\mathbf{C}(M) = \{x \in L \mid x \mathbf{C} m \quad \forall m \in M\}$  is called the *commutant* of  $M$  in  $L$  and the set  $\mathbf{Cen}(M) = \mathbf{C}(M) \cap M$  is called the *center* of  $M$ . We note that  $\mathbf{C}(L)$  is the center of  $L$  and  $\mathbf{C}(L) = \bigcap \{\mathbf{C}(a) \mid a \in L\}$ . An OML  $L$  is called *irreducible* if  $\mathbf{C}(L) = \{0, 1\}$ , and  $L$  is called *reducible* if it is not irreducible.

A *block* of an OML  $L$  is a maximal Boolean subalgebra of  $L$ . The set of all blocks of  $L$  is denoted by  $\mathfrak{A}_L$ . Note that  $\bigcup \mathfrak{A}_L = L$  and  $\bigcap \mathfrak{A}_L = \mathbf{C}(L)$ . An OML  $L$  is said to be *block-finite* if  $|\mathfrak{A}_L|$  is finite.

For any  $e$  in an OML  $L$ , the subalgebra  $S_e = [0, e'] \cup [e, 1]$  is called the (*principal*) *section generated by  $e$* . Note that for  $A, B \in \mathfrak{A}_L$ , if  $e \in A \cap B$  and  $A \cap B = S_e \cap (A \cup B)$ , then  $A \cap B = S_e \cap A = S_e \cap B$ .

**DEFINITION 1.1.** For blocks  $A, B$  of an OML  $L$  define  $A \overset{wk}{\sim} B$  if and only if  $A \cap B = S_e \cap (A \cup B)$  for some  $e \in A \cap B$ ;  $A \sim B$  if and only if  $A \neq B$  and  $A \cup B \leq L$ ;  $A \approx B$  if and only if  $A \sim B$  and  $A \cap B \neq \mathbf{C}(L)$ .

A *path* in  $L$  is a finite sequence  $B_0, B_1, \dots, B_n$  ( $n \geq 0$ ) in  $\mathfrak{A}_L$  satisfying  $B_i \sim B_{i+1}$  whenever  $0 \leq i < n$ . The path is said to *join* the blocks  $B_0$  and  $B_n$ . The number  $n$  is said to be the *length* of the path. A path is said to be *proper* if and only if  $n = 1$  or  $B_i \approx B_{i+1}$  holds whenever  $0 \leq i < n$ . A path is called to be *strictly proper* if and only if  $B_i \approx B_{i+1}$  holds whenever  $0 \leq i < n$  [1].

Let  $A, B$  be two blocks of an OML  $L$ . If  $A \sim B$  holds, then there exists a unique element  $e \in A \cap B$  satisfying  $A \cap B = (A \cup B) \cap S_e$  [1]. Using this element  $e$ , we say that  $A$  and  $B$  are *linked at  $e$*  (*strongly linked at  $e$* ) if  $A \sim B$  ( $A \approx B$ ), and use the notation  $A \sim_e B$  ( $A \approx_e B$ ). This element  $e$  is called a *vertex* of  $L$  and it is the commutator of any  $x \in A \setminus B$  and  $y \in B \setminus A$  [1]. The set of all vertices of  $L$  is denoted by  $V_L$  and  $L$  is said to be *vertex-finite* if  $|V_L|$  is finite.

Note that  $A \approx B$  implies  $A \sim B$ , and  $A \sim B$  implies  $A \overset{wk}{\sim} B$ . Some authors, for example Greechie, use the phrase “ $A$  and  $B$  meet in the section  $S_e$ ” to describe  $A \overset{wk}{\sim} B$  [3].

**DEFINITION 1.2.** Let  $L$  be an OML, and  $A, B \in \mathfrak{A}_L$ . We will say that  $A$  and  $B$  are *path-connected in  $L$* , *strictly path-connected in  $L$*  if  $A$  and  $B$  are joined by a proper path, a strictly proper path, respectively. We will say  $A$  and  $B$  are *nonpath-connected* if there is no proper path joining  $A$  and  $B$ , and  $L$  is called *nonpath-connected* if there exist two blocks

which are nonpath-connected. An OML  $L$  is called *path-connected in  $L$* , *strictly path-connected in  $L$*  if any two blocks in  $L$  are joined by a proper path, a strictly proper path, respectively. An OML  $L$  is called *relatively path-connected* iff each  $[0, x]$  is path-connected for all  $x \in L$ .

Let  $L$  be an OML, and  $A, B, C \in \mathfrak{A}_L$ . If  $A$  and  $B$  are joined with a strictly proper path  $A = B_0 \approx B_1 \approx \dots \approx B_{m-1} \approx B_m = B$  and if  $B$  and  $C$  are joined with a strictly proper path  $B = C_0 \approx C_1 \approx \dots \approx C_{n-1} \approx C_n = C$  then  $A$  and  $C$  are strictly path-connected by the *concatenated path*  $A = B_0 \approx B_1 \approx \dots \approx B_{m-1} \approx B \approx C_1 \approx \dots \approx C_{n-1} \approx C_n = C$ .

The following propositions are known.

**PROPOSITION 1.3.** *Every finite direct product of path-connected OMLs is path-connected [6].*

**PROPOSITION 1.4.** *Every infinite direct product of path-connected OMLs containing infinitely many non-Boolean factors is nonpath-connected [7].*

## 2. Finite conditions of orthomodular lattices

We present some properties of orthomodular lattices with finite conditions. A *site* is a subalgebra of an OML  $L$  of the form  $S = A \cap B$  where  $A$  and  $B$  are distinct blocks of  $L$ . An OML  $L$  is called *with uniformly finite sites* if and only if there exists a natural number  $n$  such that for all distinct blocks  $A, B$  of  $L$ ,  $|A \cap B| < n$ . We prove the following: every chain-finite OML is path-connected; every finite block of an OML  $L$  is path-connected with at least one other block in  $L$ ; every OML with uniformly finite sites is path-connected.

Greechie proved that every chain-finite OML is weakly path-connected [3]. We will prove a stronger result saying that every chain-finite OML is path-connected by using the following structure theorem of chain-finite OMLs.

**PROPOSITION 2.1.** *Every chain-finite OML  $L$  has a unique orthogonal*

decomposition  $L = L_0 \oplus L_1 \oplus \cdots \oplus L_n$  ( $0 \leq n$ ) where  $L_0$  is a Boolean algebra and  $L_1, L_2, \dots, L_n$  are simple non-Boolean chain-finite OMLs.

PROOF. Let  $L$  be a chain-finite OML. Then there exist only finitely many distinct minimal non-zero elements  $c_1, c_2, \dots, c_k \in \mathbf{C}(L)$ , and  $L = \bigoplus_{i=1}^k [0, c_i]$ . Let  $I = \{i \mid c_i \text{ is an atom of } L\}$ ,  $L_0 = \bigoplus_{i \in I} [0, c_i]$  and  $J = \{i \mid 1 \leq i \leq k, i \notin I\}$ . Then  $L = L_0 \oplus \bigoplus_{j \in J} [0, c_j]$  where  $L_0$  is a Boolean algebra since each  $[0, c_i] \subset \mathbf{C}(L)$  ( $i \in I$ ) and  $[0, c_j]$  is an irreducible non-Boolean chain-finite OML for each  $j \in J$ . Let  $n = |J|$ . This completes the proof since every irreducible chain-finite OML is simple [5].  $\square$

We prove that every chain-finite OML is path-connected using the following lemmas (2.2) and (2.3).

LEMMA 2.2. *Let  $L$  be an OML, and  $A, B \in \mathfrak{A}_L$ . If  $A \cap B = \mathbf{C}(L)$  and  $A \cup B \not\leq L$ , then there exist  $C, D \in \mathfrak{A}_L$  such that  $A \cap C \neq \mathbf{C}(L)$ ,  $C \cap D \neq \mathbf{C}(L)$  and  $D \cap B \neq \mathbf{C}(L)$ .*

PROOF. There exist  $c, d$  such that  $c, d \in A \cup B$  and  $c \vee d \notin A \cup B$  since  $A \cup B \not\leq L$ . Hence  $c \vee d \notin \mathbf{C}(L) = \bigcap \mathfrak{A}_L$ . We may assume that  $c \in A \setminus B$  and  $d \in B \setminus A$ . Therefore there exist  $C, D \in \mathfrak{A}_L$  such that  $c, c \vee d \in C$  and  $d, c \vee d \in D$ . Then  $c, d, c \vee d \notin \mathbf{C}(L)$  with  $c \in A \cap C$ ,  $c \vee d \in C \cap D$  and  $d \in D \cap B$ .  $\square$

Let  $L$  be an OML. A subalgebra  $S$  of  $L$  is said to be a *full subalgebra* if every block of  $S$  is a block of  $L$ . Note that each  $\mathbf{C}(x)$  is a full subalgebra of  $L$  for all  $x \in L$  since  $\mathfrak{A}_{\mathbf{C}(x)} = \{B \in \mathfrak{A}_L \mid x \in B\}$ .

LEMMA 2.3. *Let  $L$  be an OML. If  $[0, x]$  is path-connected  $\forall x \in L \setminus \mathbf{C}(L)$ , then  $L$  is path-connected.*

PROOF. Let  $A, B \in \mathfrak{A}_L$ . First, if  $A \cap B \neq \mathbf{C}(L)$ , then there exists  $y \in A \cap B \setminus \mathbf{C}(L)$ . Since  $y, y' \notin \mathbf{C}(L)$ ,  $[0, y]$  and  $[0, y']$  are path-connected by the hypothesis. Thus  $\mathbf{C}(y)$  is path-connected by proposition (1.3) since  $\mathbf{C}(y) = [0, y] \oplus [0, y']$ . Thus  $A$  and  $B$  are path-connected in  $\mathbf{C}(y)$  and therefore in  $L$  since  $\mathbf{C}(y)$  is a full subalgebra of  $L$ . Second, if  $A \cap B = \mathbf{C}(L)$  and  $A \cup B \leq L$ , then  $A$  and  $B$  are path-connected. Finally, if  $A \cap B = \mathbf{C}(L)$  and  $A \cup B \not\leq L$ , then there exist  $C, D \in \mathfrak{A}_L$  such that

$A \cap C \neq \mathbf{C}(L)$ ,  $C \cap D \neq \mathbf{C}(L)$  and  $D \cap B \neq \mathbf{C}(L)$  by lemma (2.2). Thus  $A$  and  $B$  are path-connected by a concatenated path by the first case.  $\square$

We are ready to prove one of our main theorems.

**THEOREM 2.4.** *Every chain-finite OML is path-connected.*

**PROOF.** Let  $L_1$  be a chain-finite OML which is nonpath-connected. So  $L_1$  is non-Boolean. We construct an infinite sequence  $x_1 > x_2 > x_3 > \dots$  by induction. There exists at least one  $x_1 \in L_1 \setminus \mathbf{C}(L_1)$  such that  $L_1[0, x_1]$  is not path-connected by lemma (2.3) since  $L_1$  is not path-connected. Let  $L_2 = L_1[0, x_1]$ . Assume that there exist  $x_i \in L_i \setminus \mathbf{Cen}(L_i)$  such that  $L_i[0, x_i]$  is not path-connected in  $L_i$  for  $1 \leq i \leq n$  and  $x_1 > x_2 > \dots > x_{n-1} > x_n$ . Let  $L_{n+1} = L_n[0, x_n]$ . Then there exists  $x_{n+1} \in L_{n+1} \setminus \mathbf{Cen}(L_{n+1})$  such that  $L_{n+1}[0, x_{n+1}]$  is not path-connected in  $L_{n+1}$  and  $x_1 > x_2 > \dots > x_n > x_{n+1}$  by lemma (2.3) since  $L_{n+1}$  is not path-connected. Thus we have an infinite sequence  $x_1 > x_2 > x_3 > \dots$  contradicting the chain-finiteness of  $L$ . This completes the proof.  $\square$

We need the following lemma to prove theorem (2.6).

**LEMMA 2.5.** *Every OML  $L$  containing a finite block has a unique orthogonal decomposition  $L = L_0 \oplus L_1 \oplus \dots \oplus L_n$  ( $0 \leq n$ ) where  $L_0$  is a Boolean algebra and  $L_1, L_2, \dots, L_n$  are irreducible OMLs each containing a finite block.*

**PROOF.** Let  $A$  be a finite block of  $L$ . Then there exist only finitely many distinct minimal non-zero elements  $c_1, c_2, \dots, c_k \in \mathbf{C}(L)$  since  $\mathbf{C}(L) \subseteq A$ , and  $L = \bigoplus_{i=1}^k [0, c_i]$ . Let  $I = \{i \mid c_i \text{ is an atom of } L\}$ ,  $L_0 = \bigoplus_{i \in I} [0, c_i]$  and  $J = \{i \mid 1 \leq i \leq k, i \notin I\}$ . Then  $L = L_0 \oplus \bigoplus_{j \in J} [0, c_j]$  where  $L_0$  is a Boolean algebra since each  $[0, c_i] \subset \mathbf{C}(L)$  ( $i \in I$ ) and  $[0, c_j]$  is an irreducible OML containing a finite block. Let  $n = |J|$ . This completes the proof.  $\square$

**THEOREM 2.6.** *Each finite block of a non-Boolean OML  $L$  is path-connected with at least one other block of  $L$ .*

PROOF. We may assume that  $|L| \geq 4$ , and that  $L$  is an irreducible OML containing a finite block  $A$  by proposition (1.3) and lemma (2.5). We will prove the claim by induction on the cardinality of  $A$ . If  $|A| = 2^2$ , then  $A$  is an horizontal summand of  $L$  so that  $A$  is path-connected with each block in  $L$ . Assume every finite block  $A$  with  $|A| \leq 2^n$  ( $n \geq 2$ ) of any non-Boolean OML is path-connected with at least one other block of that OML. Let  $|A| = 2^{n+1}$ . Let us show that  $A$  is path-connected with at least one other block of  $L$ . First, if  $A \cap B = \{0, 1\}$  for each  $B \in \mathfrak{A}_L \setminus \{A\}$ , then we claim that  $\bigcup \mathfrak{A}_L \setminus \{A\}$  is a subalgebra of  $L$ . Let  $M = \bigcup \mathfrak{A}_L \setminus \{A\}$ . Suppose  $M$  is not a subalgebra of  $L$ , then there exist two distinct element  $a, b$  in  $M$  such that  $a \vee b \notin M$  and  $a \vee b \in A$ . Since  $0 < a < a \vee b < 1$ , there exist a block  $C$  containing  $a$  and  $a \vee b$ . Thus  $A \neq C$  since  $a \in C \setminus A$ , and  $C \subset M$ . Therefore  $a \vee b \in A \cap C \setminus \{0, 1\}$ . This contradicts the assumption  $A \cap B = \{0, 1\}$  for each  $B \in \mathfrak{A}_L \setminus \{A\}$ . Thus  $A$  is a horizontal summand of  $L$ . Therefore we may assume  $A \cap B \neq \{0, 1\}$  for some  $B \in \mathfrak{A}_L \setminus \{A\}$ . Let  $x \in A \cap B \setminus \{0, 1\}$ . Then  $\mathbf{C}(x) = [0, x] \oplus [0, x']$  and  $A = A[0, x] \oplus A[0, x']$ . If  $[0, x]$  is Boolean, then  $[0, x']$  is path-connected by induction hypothesis since  $|[0, x']| \leq 2^n$ . Thus  $\mathbf{C}(x) = [0, x] \oplus [0, x']$  is path-connected by proposition (1.3). Therefore we may assume that  $[0, x]$  is non-Boolean. Then  $A[0, x]$  is path-connected with another block  $D \in \mathfrak{A}_{\mathbf{C}(x)[0, x]}$  by the induction hypothesis since  $|A[0, x]| \leq 2^n$ . Therefore  $A = A[0, x] \oplus A[0, x']$  is path-connected with  $D \oplus A[0, x'] \in \mathfrak{A}_{\mathbf{C}(x)}$  in  $\mathbf{C}(x)$  by proposition (1.3). Thus  $A = A[0, x] \oplus A[0, x']$  is path-connected with  $D \oplus A[0, x']$  in  $L$  since  $\mathbf{C}(x)$  is a full subalgebra of  $L$ . This completes the proof.  $\square$

We have the following class of path-connected OMLs with a finite condition.

**THEOREM 2.7.** *Every OML  $L$  with uniformly finite sites is path-connected.*

PROOF. We may assume that  $|\mathfrak{A}_L| \geq 2$ . We will prove the claim by induction on the maximum cardinality  $2^n$  of  $|A \cap B|$  for all distinct blocks  $A, B$  of  $L$ . Let  $L$  be an OML such that  $|A \cap B| \leq 2^n$  for all distinct blocks  $A, B$  in  $L$ . If  $n = 1$ , then  $L$  is path-connected since  $L$  is a horizontal sum of Boolean algebras. Assume  $n > 1$ , and assume that every OML such that  $|A \cap B| \leq 2^{n-1}$  for all distinct blocks  $A, B$  of that OML is path-connected. If  $L$  is an OML such that  $|A \cap B| \leq 2^n$  for all

distinct blocks  $A, B$  of  $L$ . First, assume  $A, B$  to be two distinct blocks of  $L$  such that  $A \cap B \neq \mathbf{C}(L)$ . If  $A \cup B \leq L$ , then  $A$  and  $B$  are path-connected. Thus we may assume that  $A \cup B \not\leq L$ . Let  $x \in (A \cap B) \setminus \mathbf{C}(L)$ . Then  $[0, x]$  and  $[0, x']$  are path-connected by inductive hypothesis since  $|C \cap D| \leq 2^{n-1}$  for all distinct blocks  $C, D$  of  $[0, x]$  and  $|E \cap F| \leq 2^{n-1}$  for all distinct blocks  $E, F$  of  $[0, x']$ . Thus  $\mathbf{C}(x) = [0, x] \oplus [0, x']$  is path-connected by proposition (1.4). Hence  $A$  and  $B$  are path-connected in  $\mathbf{C}(x)$ , and therefore in  $L$  since  $\mathbf{C}(x)$  is a full subalgebra of  $L$ . Finally, assume  $A \cap B = \mathbf{C}(L)$ . If  $A \cup B \leq L$ , then  $A$  and  $B$  are path-connected. Now we may assume that  $A \cup B \not\leq L$  and  $A \cap B = \mathbf{C}(L)$ . Then there exist  $G, H \in \mathfrak{A}_L$  such that  $A \cap G \neq \mathbf{C}(L)$ ,  $G \cap H \neq \mathbf{C}(L)$ , and  $H \cap B \neq \mathbf{C}(L)$  by lemma (2.2). Thus  $A$  and  $B$  are path-connected by a concatenated path by the first case. This completes the proof.  $\square$

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