KYUNGPOOK Math. J. 57(2017), 145-161 https://doi.org/10.5666/KMJ.2017.57.1.145 pISSN 1225-6951 eISSN 0454-8124 © Kyungpook Mathematical Journal

## Polynomials and Homotopy of Virtual Knot Diagrams

Myeong-Ju Jeong

Department of Mathematics and Computer Science, Korea Science Academy of KAIST 111 Baekyang Gwanmun-Ro, Busanjin-Gu, Busan 614-822 Korea e-mail: mjjeong@kaist.ac.kr

Chan-Young Park\*

Department of Mathematics, College of Natural Sciences Kyungpook National University, Daegu 702-701 Korea e-mail: chnypark@knu.ac.kr

Maeng Sang Park

Department of Mathematics Pusan National University, Pusan 609-735 Korea e-mail: msjocund@hanmail.net

ABSTRACT. If a virtual knot diagram can be transformed to another virtual one by a finite sequence of crossing changes, Reidemeister moves and virtual moves then the two virtual knot diagrams are said to be *homotopic*. There are infinitely many homotopy classes of virtual knot diagrams.

We give necessary conditions by using polynomial invariants of virtual knots for two virtual knots to be homotopic. For a sequence S of crossing changes, Reidemeister moves and virtual moves between two homotopic virtual knot diagrams, we give a lower bound for the number of crossing changes in S by using the affine index polynomial introduced in [13]. In [10], the first author gave the q-polynomial of a virtual knot diagram to find Reidemeister moves of virtually isotopic virtual knot diagrams. We find how to apply Reidemeister moves by using the q-polynomial to show homotopy of two virtual knot diagrams.

#### 1. Introduction

Any classical knot diagram can be unknotted by a sequence of Reidemeister moves and crossing changes. So crossing change is called an unknotting operation

Received August 25, 2015; revised January 29, 2016; accepted May 4, 2016. 2010 Mathematics Subject Classification: 57M25, 57M27.

Key words and phrases: affine index polynomial, virtual homotopy, crossing change, Gordian distance.

<sup>\*</sup> Corresponding Author.

for classical knot diagrams. There was an open question whether one can get the minimal number of crossing changes from a knot diagram with a minimal crossing number to get the trivial knot diagram. For this question S. Bleiler [1] and Y. Nakanishi [19] gave a counterexample C(5,1,4) in Conway notation in Figure 1. Crossing change is no more an unknotting operation for virtual knot diagrams. Two virtual knot diagrams are homotopic if there is a sequence of Reidemeister moves, virtual moves and crossing changes. We will give necessary conditions for two given virtual knot diagrams to be homotopic. To show homotopy of virtual knot diagrams, we investigate how to apply crossing changes by using an affine index polynomial  $P_K(t)$  of a virtual knot diagram K and how to apply Reidemeister moves by using a polynomial  $Q_K(t)$  of a virtual knot diagram K.

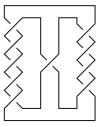


Figure 1: A knot C(5, 1, 4).

In 1996 L. H. Kauffman introduced virtual knots which generalize classical knots [12]. A *virtual knot diagram* is a knot diagram allowed to have virtual crossings. We denote a virtual crossing by an encircled crossing as shown in Figure 2.

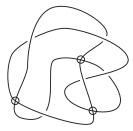


Figure 2: A virtual knot diagram.

If a knot K is isotopic to another knot K' then there is a sequence of moves from a diagram of K to a diagram of K' as shown in Figure 3. These moves are called *Reidemeister moves*. The moves of diagrams shown in Figure 4 are called *virtual moves*. Two virtual knot diagrams  $K_1$  and  $K_2$  are said to be *virtually isotopic* if there is a sequence of Reidemeister moves and virtual moves from  $K_1$  to  $K_2$ . A *virtual knot* is defined to be the virtual isotopy class of a virtual knot diagram.

From now on all virtual knot diagrams are assumed to be oriented. We define

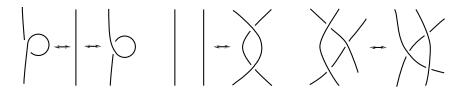


Figure 3: Reidemeister moves.

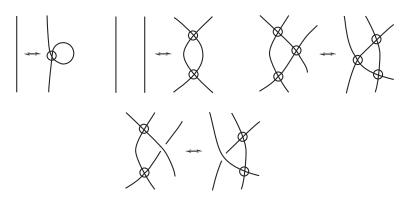


Figure 4: Virtual moves.

the  $sign \operatorname{sgn}(c)$  of a crossing c of a virtual knot diagram as shown in Figure 5. The  $writhe\ w(K)$  of a virtual knot diagram K is defined to be the sum of signs of all crossings of K.



Figure 5: The sign of a crossing.

M. Goussarov, M. Polyak and O. Viro defined finite type invariants of virtual knots and gave some combinatorial representations of finite type invariants of low degree by using Gauss diagrams [6]. The  $Gauss\ diagram$  of a virtual knot diagram K is an oriented circle with chords corresponding to crossings. The two endpoints of a chord correspond to the preimages of the crossing of K. A chord corresponding to a crossing c is oriented from the preimage of the over crossing point of c to the preimage of the under crossing point of c. A chord is assumed to have the sign of the crossing corresponding to the chord. See Figure 6. We denote the Gauss

diagram of K by G(K).

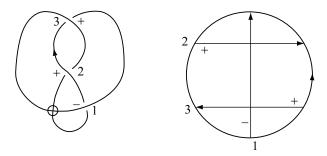


Figure 6: A Gauss diagram of a virtual knot diagram.

A virtual knot can be represented as a Gauss diagram and vice versa [6, 12]. For a Gauss diagram G, there are many virtual knot diagrams whose Gauss diagrams are the same G. All these virtual knot diagrams are virtually isotopic [6]. In Gauss diagrams, the Reidemeister moves can be described as a sequence of moves shown in Figure 7 [6, 21]. In the figure  $\epsilon$  denotes the sign of a chord which can be either + or -.

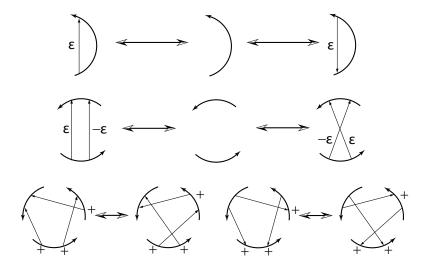


Figure 7: Moves of Gauss diagrams.

A local move of virtual knot diagrams shown in Figure 8 is called a *crossing change* or a CC-move. Two virtual knot diagrams are said to be *homotopic* if they are related by a sequence of CC-moves, Reidemeister moves, and virtual moves. If two virtual knot diagrams  $K_1$  and  $K_2$  are homotopic then the *Gordian distance* 



Figure 8: A crossing change.

 $d_G(K_1, K_2)$  between  $K_1$  and  $K_2$  is defined to be the minimal number of CC-moves needed to transform  $K_1$  to  $K_2$ , up to Reidemeister moves and virtual moves. Although any two classical knot diagrams are homotopic, we will see that there are infinitely many non-homotopic pairs of virtual knot diagrams.

A. Kawauchi gave a condition on a pair of the Alexander polynomials of knots which are realizable by a pair of knots with Gordian distance one [15]. He showed that there are infinitely many mutually disjoint infinite subsets of the set of the Alexander polynomials of knots such that every pair of distinct polynomials in each subset is not realizable by any pair of knots with Gordian distance one [15]. Y. Miyazawa gave several evaluations of the Gordian distance of two classical knots by using the HOMFLY polynomial, the Jones polynomial and the Q-polynomial [18]. Similarly to the classical case, we can extend virtual knot diagrams to virtual link diagrams with multiple components naturally [12]. Two virtual link diagrams are said to be welded equivalent if one can be transformed to the other by a sequence of Reidemeister moves, virtual moves and the upper forbidden move shown in Figure 9. Two welded link diagrams are said to be homotopic if one can be transformed to the other by a sequence of Reidemeister moves, virtual moves and self crossing changes. H. A. Dye and L. H. Kauffman extended Milnor's  $\mu$  and  $\bar{\mu}$  invariants to welded and virtual links [5]. Both the upper forbidden move and the lower forbidden move in Figure 9 are called forbidden moves which unknot all virtual knots [8, 14, 20]. A  $\Delta$ move is a local move of virtual knot diagrams as shown in Figure 10. T. Kanenobu showed that a  $\Delta$ -move can be realized by a finite sequence of the Reidemeister moves, the virtual moves and the forbidden moves [14]. The first author showed that the values of a Vassiliev invariant of degree 2 for two virtual knots related by a  $\Delta$ -move differ by 48 [9].

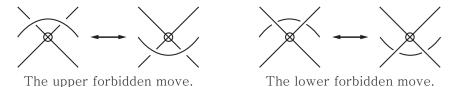


Figure 9: Forbidden moves.

A. Henrich defined a polynomial invariant of virtual knots by using weight of a crossing. Z. Cheng [2] introduced an odd writhe polynomial  $f_K(t)$  of virtual knots

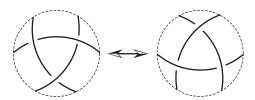


Figure 10: A  $\Delta$ -move.

K considering the Gauss diagram of a virtual knot and coloring of arcs. Z. Cheng and H. Gao [3] generalized the odd writhe polynomial to a sequence of polynomials from which Henrich's polynomial can be recovered. L. H. Kauffman introduced the affine index polynomial of a virtual knot based on a Cheng coloring of a virtual knot diagram [13]. An affine index polynomial  $P_K(t)$  is given in the form

$$P_K(t) = \sum_{c} \text{sgn}(c) (t^{W_K(c)} - 1),$$

where  $\operatorname{sgn}(c)$  and  $W_K(c)$  are the sign and the weight of a crossing c respectively. We will show that  $P_{K_2}(t) - P_{K_1}(t) = \pm (t^n + t^{-n} - 2)$  for some integer n if  $K_2$  is obtained from  $K_1$  by applying a CC-move.

The first author defined a polynomial  $q_K(t)$  of a virtual knot diagram K in [10]. If  $K_1$  and  $K_2$  are related by a second Reidemeister move then  $q_{K_2}(t) - q_{K_1}(t) = \pm 1$  or 0. If  $K_1$  and  $K_2$  are related by a third Reidemeister move then  $q_{K_2}(t) - q_{K_1}(t) = \pm (t^{\alpha+\beta} - t^{\alpha} - t^{\beta})$  for some  $\alpha, \beta \in \mathbb{Z}$ . To show that two virtual knot diagrams  $K_1$  and  $K_2$  are homotopic, basically we find a sequence of CC-moves, Reidemeister moves and virtual moves between  $K_1$  and  $K_2$ . The affine index polynomial and q-polynomial can shed light on finding such a sequence.

In Section 2, we give a sequence of virtual knots every pair of which is non-homotopic and give necessary condition for two virtual knot diagrams to be homotopic by using the affine index polynomial. Moreover we give a lower bound for the number of CC-moves needed to transform a virtual knot diagram to a homotopic one by comparing their affine index polynomials. In Section 3, we give a relationship between  $q_{K_1}(t)$  and  $q_{K_2}(t)$  if  $K_1$  and  $K_2$  differ by a CC-move. When two virtual knot diagrams  $K_1$  and  $K_2$  are homotopic, their Gauss diagrams are related by the moves corresponding to CC-moves and Reidemeister moves.  $q_K(t)$  is defined from the Gauss diagram of K and it is sensitive to Reidemeister moves. So we can predict how virtual knot diagrams  $K_1$  and  $K_2$  are related by Reidemeister moves by comparing  $q_{K_1}(t)$  and  $q_{K_2}(t)$  if  $K_1$  and  $K_2$  are virtually isotopic. We can map out a sequence of Reidemeister moves and CC-moves from  $K_1$  to  $K_2$  by using the affine index polynomial and the q-polynomial for homotopic virtual knot diagrams  $K_1$  and  $K_2$ . We give an example illustrating that  $P_K(t)$  and  $q_K(t)$  are useful to show that two virtual knot diagrams are homotopic.

## 2. Polynomial Invariants and CC-moves of Virtual Knot Diagrams

Y. Miyazawa found a 2-variable polynomial invariant of virtual links [16] and generalized it by constructing a multi-variable polynomial invariant of virtual links [17]. He gave a lower bound on the virtual crossing number by using the multi-variable polynomial.

H. A. Dye and L. H. Kauffman defined the arrow polynomial of a virtual link which is a generalization of the bracket polynomial [4]. By normalizing the arrow polynomial we get an invariant of virtual knots and links. By changing variables suitably, we may get the normalized arrow polynomial from the Miyazawa's multivariable polynomial and vice versa. By using the normalized arrow polynomial we will get a necessary condition for two given virtual knot diagrams to be homotopic.

**Definition 2.1.** The *arrow polynomial* of a virtual link diagram is defined by using the following relations.

- (1)  $\langle L_{+} \rangle = A \langle L_{0} \rangle + A^{-1} \langle L_{\infty} \rangle$  and  $\langle L_{-} \rangle = A^{-1} \langle L_{0} \rangle + A \langle L_{\infty} \rangle$ , where  $L_{+}$ ,  $L_{-}$ ,  $L_{0}$ , and  $L_{\infty}$  are virtual link diagrams which differ as shown in Figure 11 and A is an indeterminate.
- (2)  $\langle C_1 \rangle = \langle C_2 \rangle$  and  $\langle C_3 \rangle = \langle C_4 \rangle$ , where  $C_1$ ,  $C_2$ ,  $C_3$  and  $C_4$  are virtual link diagrams which differ as shown in Figure 12.
- (3)  $\langle O \rangle = 1$  and  $\langle L \cup O_m \rangle = (-A^2 A^{-2})X_m \langle L \rangle$ , where  $X_m$ 's are indeterminates, O is the trivial knot diagram with no crossings and no poles and  $O_m$  is the diagram with 2m poles as shown in Figure 12.

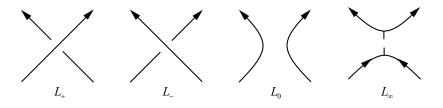


Figure 11:

The arrow polynomial is invariant under Reidemeister moves and virtual moves except for the first Reidemeister moves [4]. Similarly to the bracket polynomial we can normalize the arrow polynomial to get an invariant of virtual links. The normalized arrow polynomial  $\langle L \rangle_{NA}$  is defined by the formula

$$\langle L \rangle_{NA} = (-A^3)^{-w(L)} \langle L \rangle.$$

The arrow polynomial and normalized arrow polynomial take values in the polynomial ring  $\mathbb{Z}[A, A^{-1}, X_1, X_2, \cdots]$ .

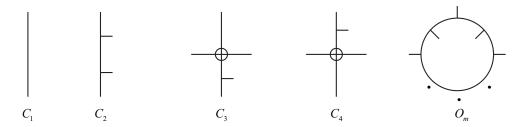


Figure 12:

The first and second authors defined polar link diagrams, a generalization of virtual link diagrams, and extended the normalized arrow polynomial to polar link diagrams. We gave the following lemma, which is useful to see whether two virtual knot diagrams are homotopic or not [11].

**Lemma 2.2.**([11]) If L and L' are homotopic polar link diagrams then

$$\langle L \rangle_{NA} \equiv \langle L' \rangle_{NA} \mod A^4 - 1.$$

In particular, if two virtual knot diagrams K and K' are homotopic then their normalized arrow polynomials  $\langle K \rangle_{NA}$  and  $\langle K' \rangle_{NA}$  are congruent modulo  $A^4 - 1$ . In the following example, we show that there are infinitely many homotopy classes of virtual knot diagrams.

**Example 2.3.** Consider a sequence of virtual knot diagrams  $\{K_n\}_{n=0}^{\infty}$  as shown in Figure 13. Inductively we can show that the degree of  $X_1$  of  $\langle K_n \rangle_{NA}$  modulo  $A^4 - 1$  is n. Therefore any two virtual knot diagrams in the sequence are not homotopic by Lemma 2.2.

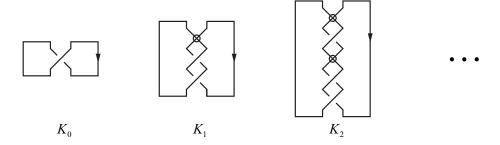


Figure 13: A sequence of virtual knot diagrams with distinct homotopy classes.

A *flat crossing* is an intersecting line segments with no information about over and under strand. A *flat virtual knot diagram* is a virtual knot diagram whose

crossings are replaced with flat crossings. From a virtual knot diagram K we get a flat virtual knot diagram F(K) by taking its shadow of K. Figure 14 shows a flat virtual knot diagram with two flat crossings.

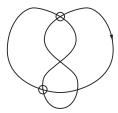


Figure 14: A flat virtual knot.

By replacing crossings in Reidemeister moves and virtual moves with flat crossings, we get *flat Reidemeister moves* and *flat virtual moves* respectively. A sequence of flat Reidemeister moves and flat virtual moves is called a *flat virtual isotopy*. A *flat virtual knot* is defined as the flat virtual isotopy class of a flat virtual knot diagram.

Let K be a virtual knot diagram. Then F(K) can be regarded as a 4-regular graph G by considering a flat crossing as a vertex. An arc of F(K) is an edge of G that it represents. Label each arc of F(K) with an integer so that it satisfies the rule in Figure 15. L. H. Kauffman proved that such labeling exists always [13]. Such labeling of arcs for a flat virtual knot diagram is called a *Cheng coloring*. See Figure 16.

Let c be a flat crossing of a labeled flat virtual knot diagram. If arcs near c is labeled as shown in the left part of Figure 15, then we define  $W_+(c) = m - n - 1$  and  $W_-(c) = n - m + 1$ .  $W_+(c)$  and  $W_-(c)$  are independent of Cheng colorings of F(K) [3, 13].

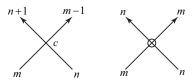


Figure 15: Labeling of arcs.

For each crossing c of K, there is a corresponding flat crossing c' in F(K). If F(K) is labeled as previously described, then define  $W_K(c) = W_{\operatorname{sgn}(c)}(c')$ . For simplicity sake we often denote  $W_K(c)$  by W(c). We denote the set of all crossings of K by C(K). The affine index polynomial  $P_K(t)$  of K is defined by the equation

$$P_K(t) = \sum_{c \in C(K)} \operatorname{sgn}(c) (t^{W_K(c)} - 1).$$

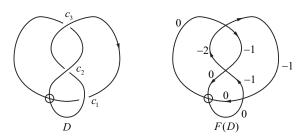


Figure 16: A Cheng coloring.

L. H. Kauffman showed that  $P_K(t)$  is an invariant of virtual knots [13]. For example the affine index polynomial of a virtual knot diagram D in Figure 16 can be calculated as following

$$P_D(t) = \operatorname{sgn}(c_1)(t^{W_K(c_1)} - 1) + \operatorname{sgn}(c_2)(t^{W_K(c_2)} - 1) + \operatorname{sgn}(c_3)(t^{W_K(c_3)} - 1)$$
  
=  $-t - t^{-1} + 2$ .

We give a necessary condition for two given virtual knot diagrams  $K_1$  and  $K_2$  to be homotopic by using the affine index polynomials of  $K_1$  and  $K_2$ .

**Lemma 2.4.** Let  $K_1$  be a virtual knot diagram and  $c_1 \in C(K_1)$ . If  $K_2$  is the virtual knot diagram obtained from  $K_1$  by changing the crossing  $c_1$  then

$$P_{K_2}(t) - P_{K_1}(t) = -\operatorname{sgn}(c_1)(t^{W_{\operatorname{sgn}(c_1)}(c_1)} + t^{-W_{\operatorname{sgn}(c_1)}(c_1)} - 2).$$

*Proof.* Assume that a crossing  $c_1$  of  $K_1$  is changed to a crossing  $c_2$  of  $K_2$  by the CC-move. Since  $F(K_1) = F(K_2)$ ,  $\operatorname{sgn}(c_1) = -\operatorname{sgn}(c_2)$  and  $W_+(c) = -W_-(c)$ , we see that

$$\begin{array}{lcl} P_{K_2}(t) - P_{K_1}(t) & = & \mathrm{sgn}(c_2)(t^{W_{\mathrm{sgn}(c_2)}(c_2)} - 1) - \mathrm{sgn}(c_1)(t^{W_{\mathrm{sgn}(c_1)}(c_1)} - 1) \\ & = & -\mathrm{sgn}(c_1)(t^{W_{\mathrm{sgn}(c_1)}(c_1)} + t^{-W_{\mathrm{sgn}(c_1)}(c_1)} - 2). \end{array} \quad \Box$$

Now we see that  $P_{K_2}(t) - P_{K_1}(t) = \pm (t^{-n} + t^n - 2)$  for some integer n if two virtual knot diagrams  $K_1$  and  $K_2$  are related by a single CC-move. Moreover Either n or -n is the weight of the crossing which was changed by the CC-move. Therefore we get the following

**Theorem 2.5.** Let  $K_1$  and  $K_2$  be homotopic virtual knot diagrams. Then

$$P_{K_1}(t) - P_{K_1}(t^{-1}) = P_{K_2}(t) - P_{K_2}(t^{-1}).$$

Moreover

$$d_G(K_1, K_2) \ge \sum_{i=1}^m |a_i|,$$

where  $P_{K_2}(t) - P_{K_1}(t) = \sum_{i=1}^m a_i(t^{-i} + t^i - 2)$ .

**Example 2.6.** Let  $\{K_n\}_{n=0}^{\infty}$  be a sequence of virtual knot diagrams as shown in Figure 17. We can easily see that  $P_{K_n}(t) = -n(t^{-1} + t^1 - 2)$  for each  $n = 0, 1, \cdots$ . Therefore  $d_G(K_n, O) \geq n$  by Theorem 2.5, where O is the trivial virtual knot diagram. Since  $K_n$  can be transformed to O by changing n crossings, we see that  $d_G(K_n, O) = n$ .

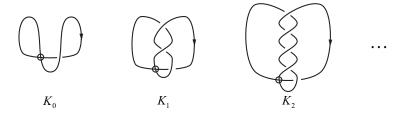


Figure 17:

### 3. CC-moves and Reidemeister Moves

If two virtual knot diagrams  $K_1$  and  $K_2$  are supposed to be homotopic, then we usually want to find a sequence of Reidemeister moves, virtual moves and CC-moves. This can be done through their Gauss diagrams. In Gauss diagrams we need not to consider virtual moves by its definition. So we are interested in finding a sequence of Reidemeister moves and CC-moves in a Gauss diagrammatic approach. The affine index polynomial gives us information on CC-moves but it does not give us any information about Reidemeister moves as it is an invariant. In this section we find how to apply Reidemeister moves to show homotopy of  $K_1$  and  $K_2$  by using q-polynomials.

A q-polynomial  $q_K(t)$  of a virtual knot diagram K is defined to study Reidemeister moves [10]. It is invariant under the first Reidemeister move but not invariant under the second Reidemeister move and the third Reidemeister move. It is defined by using mixed pairs of crossings of a virtual knot diagram and degrees of its crossings. In [10], the first author defined the degree  $\deg(c)$  of a crossing c by using Gauss diagram. A Cheng coloring of a virtual knot diagram K can be represented in a Gauss diagram [3] and we can see that  $\deg(c) = -W(c)$  for any crossing c. Refer [3, 10] for more details.

Let K be a virtual knot diagram with crossings  $c_1$  and  $c_2$ . If the two chords corresponding to  $c_1$  and  $c_2$  appear as shown in Figure 18, then the ordered pair  $(c_1, c_2)$  is called a *mixed pair of crossings*. For example the knot in Figure 6 has two mixed pairs  $(c_1, c_3)$ ,  $(c_2, c_1)$  of crossings.

Denote the set of all mixed pairs of crossings of K by M(K). The two chords of a Gauss diagram G(K) corresponding to crossings  $c_1$  and  $c_2$  of K intersect if  $(c_1, c_2) \in M(K)$ .



Figure 18:

Let c be a crossing of a virtual knot diagram K. For simplicity sake we will denote the chord of the Gauss diagram G(K) corresponding to the crossing c by c. For a chord c of a Gauss diagram, assume that we walk on the chord c from the starting point of c to the ending point of c. Let  $r_+$  and  $r_-$  be the numbers of positive chords and negative chords from left to right respectively and let  $l_+$  and  $l_-$  be the numbers of positive chords and negative chords from right to left respectively. We define the  $degree \deg(c)$  of a crossing c by the equation

$$\deg(c) = r_{+} - r_{-} - l_{+} + l_{-}.$$

For  $c_1, c_2 \in C(K)$  we also define  $\deg(c_1, c_2)$  by the formula

$$\deg(c_1, c_2) = \deg(c_1) - \deg(c_2).$$

For example, if  $c_1, c_2, c_3$  and  $c_4$  are crossings of a virtual knot diagram whose Gauss diagrams are as shown in Figure 19, then  $\deg(c_1) = 1$ ,  $\deg(c_2) = 0$ ,  $\deg(c_3) = -2$ ,  $\deg(c_4) = 1$ ,  $\deg(c_1, c_2) = 1$ ,  $\deg(c_2, c_3) = 2$  and  $\deg(c_3, c_4) = -3$ .

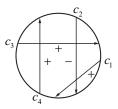


Figure 19:

Now we define a polynomial  $q_K(t)$  associated to K by the equation

$$q_K(t) = \sum_{(c_1, c_2) \in M(K)} \operatorname{sgn}(c_1) \operatorname{sgn}(c_2) t^{\deg(c_1, c_2)}.$$

**Lemma 3.1.**([10]) Let  $K_1$  and  $K_2$  be related by a second Reidemeister move and their Gauss diagrams be as shown in Figure 20. Then

$$q_{K_1}(t) - q_{K_2}(t) = -1.$$

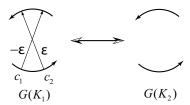


Figure 20:

**Lemma 3.2.**([10]) Let  $K_1$  and  $K_2$  be related by a second Reidemeister move and their Gauss diagrams be as shown in Figure 21. Then

$$q_{K_1}(t) = q_{K_2}(t).$$

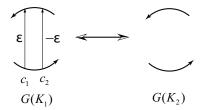


Figure 21:

**Lemma 3.3.**([10]) Let two virtual knot diagrams  $K_1$  and  $K_2$  be related by a third Reidemeister move. Then there exist  $\alpha, \beta \in \mathbb{Z}$  such that

$$q_{K_1}(t) - q_{K_2}(t) = \pm (t^{\alpha+\beta} - t^{\alpha} - t^{\beta}).$$

Let  $K_1$  and  $K_2$  be virtual knot diagrams related by a third Reidemeister move and  $c_1, c_2, c_3$  be the crossings of  $K_1$  involved with the third Reidemeister move. Then  $\alpha$ ,  $\beta$  and  $\alpha + \beta$  are given in the form  $\deg(c_i) - \deg(c_j)$  for some distinct elements i, j of the set  $\{1, 2, 3\}$  [10].

Let  $K_1$  and  $K_2$  be virtually isotopic virtual knot diagrams. Then there exists a sequence S of Reidemeister moves transforming  $K_1$  to  $K_2$ . We denote the numbers of the first Reidemeister moves, the second Reidemeister moves and the third Reidemeister moves in the sequence S by  $n_1(S)$ ,  $n_2(S)$  and  $n_3(S)$  respectively. A lower bound for  $n_1(S)$  is the difference  $|w(K_1) - w(K_2)|$  of writhes of  $K_1$  and  $K_2$ . By using Lemma 3.1, Lemma 3.2 and Lemma 3.3, we get the following theorems for the numbers of the second Reidemeister moves and the third Reidemeister moves in S.

**Theorem 3.4.**([10]) Let  $K_1$  and  $K_2$  be two virtually isotopic virtual knot diagrams and S be a sequence of Reidemeister moves deforming  $K_1$  to  $K_2$ . Let  $g_K(t) = q_K(t) - q_K(1)$  and  $g_{K_1}(t) - g_{K_2}(t) = \sum_{i=m}^l a_i(t^i - 1)$ . Then we get inequalities

$$n_3(S) \ge \frac{\sum_{i=m}^{l} |a_i|}{3} \text{ and } n_3(S) \ge \left| \sum_{i=m}^{l} a_i \right|.$$

**Theorem 3.5.**([10]) Let  $K_1$  and  $K_2$  be two virtually isotopic knot diagrams and S be a sequence of Reidemeister moves deforming  $K_1$  to  $K_2$ . Let  $q_{K_1}(t) - q_{K_2}(t) = \sum_{i=m}^{l} a_i t^i$ . Then we get an inequality

$$n_2(S) + 3n_3(S) \ge \sum_{i=m}^{l} |a_i|.$$

There are four versions of second Reidemeister moves for oriented virtual knot diagrams as shown in Figure 22. For the numbers of second Reidemeister moves and third Reidemeister moves, C. Hayashi introduced cowrithe of a knot which can be given as  $q_K(1)$  and gave the following theorem.

**Theorem 3.6.**([7]) Let K be a virtual knot diagram. The first Reidemeister move does not change  $q_K(1)$ . The second Reidemeister moves  $\Omega 2a$  and  $\Omega 2b$  change  $q_K(1)$  by  $\pm 1$ . The second Reidemeister moves  $\Omega 2c$  and  $\Omega 2d$  do not change  $q_K(1)$ . The third Reidemeister move changes  $q_K(1)$  by  $\pm 1$ .

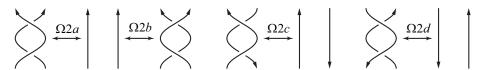


Figure 22: The second Reidemeister moves.

Let  $K_1$  and  $K_2$  be related by a CC-move. Assume that a crossing  $c_1 \in C(K_1)$  is changed to  $c_2 \in C(K_2)$  by the CC-move. Let

$$q_{K_1}(t) = \sum_{(c,c') \in M(K_1)} \operatorname{sgn}(c)\operatorname{sgn}(c')t^{\deg(c,c')} = g(t) + t^{\deg(c_1)}h(t) + t^{-\deg(c_1)}i(t),$$

where g(t), h(t) and i(t) come from the sums for  $c, c' \neq c_1$ , for  $c = c_1$  and for  $c' = c_1$  respectively. Note that if  $(c_1, c') \in M(K_1)$  then  $(c', c_2) \in M(K_2)$ . Similarly, if  $(c, c_1) \in M(K_1)$  then  $(c_2, c) \in M(K_2)$ . Since  $\operatorname{sgn}(c_2) = -\operatorname{sgn}(c_1)$  we have

$$q_{K_2}(t) = g(t) - t^{\deg(c_1)}h(t^{-1}) - t^{-\deg(c_1)}i(t^{-1}).$$

Now we get the following

**Lemma 3.7.** Let  $K_2$  be a virtual knot diagram obtained from  $K_1$  by changing a crossing  $c_1$  of  $K_1$ . Let  $h(t) = t^{-\deg(c_1)} \sum_{(c_1,c') \in M(K_1)} t^{\deg(c_1,c')}$  and  $i(t) = t^{\deg(c_1)} \sum_{(c,c_1) \in M(K_1)} t^{\deg(c,c_1)}$ . Then

$$q_{K_2}(t) - q_{K_1}(t) = -t^{\deg(c_1)}(h(t) + h(t^{-1})) - t^{-\deg(c_1)}(i(t) + i(t^{-1})).$$

Note that the coefficients of  $t^k$  and  $t^{-k}$  are the same for  $h(t) + h(t^{-1})$  and  $i(t) + i(t^{-1})$  for all  $k = 1, 2, \cdots$ .

**Example 3.8.** Let  $K_1$  and  $K_2$  be the virtual knot diagrams as shown in Figure 23. Since  $P_{K_1}(t) = t^{-2} + t^{-1} - 4 + t + t^2$  and  $P_{K_2}(t) = t^{-1} - 2 + t$ , we have

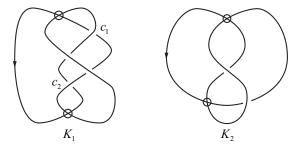


Figure 23:

 $P_{K_2}(t) - P_{K_1}(t) = -t^{-2} + 2 - t^2$ . If  $K_1$  and  $K_2$  are homotopic then we need at least 1 CC-move which changes a crossing of  $K_1$  with weight  $\pm 2$  and with sign + by Theorem 2.5. There are two crossings  $c_1$  and  $c_2$  in  $K_1$  with weight  $\pm 2$  and with sign +. If we change the crossing  $c_1$  then we get a diagram  $K_3$  in Figure 24. Since  $q_{K_3}(t) = t^{-3} - t^{-1} - 1$  and  $q_{K_2}(t) = t^{-2}$ , we have  $q_{K_2}(t) - q_{K_3}(t) = -t^{-3} + t^{-2} + t^{-1} + 1$ . By Theorem 3.4 we see that we need a third Reidemeister move at least

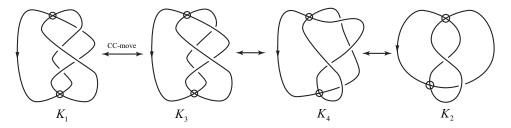


Figure 24:

once. By Theorem 3.5, we also see that we need another third Reidemeister move or a second Reidemeister move. By comparing  $w(K_3)$  and  $w(K_2)$ , we see that we

may not need a first Reidemeister move. By applying a third Reidemeister move to  $K_3$  we get  $K_4$  as shown in Figure 24. Now  $q_{K_4}(t) = t^{-2} - 1$  and  $q_{K_2}(t) - q_{K_4}(t) = 1$ . From Lemma 3.1, we see that  $K_4$  may be transformed to  $K_2$  by a suitable second Reidemeister move. Actually we get  $K_1$  by applying a second Reidemeister move to  $K_4$  as shown in Figure 24.

**Acknowledgements.** The first named author was supported by the Ministry of Science, ICT and Future Planning.

# References

- S. Bleiler, A note on unknotting number, Math. Proc. Cambridge Philos. Soc., 96(1984), 469–471.
- [2] Z. Cheng, A polynomial invariant of virtual knots, Proc. Amer. Math. Soc., 142(2)(2014), 713–725.
- [3] Z. Cheng and H. Gao, A polynomial invariant of virtual links, J. Knot Theory Ramifications, **22(12)**(2013), 1341002, 33 pp.
- [4] H. A. Dye and L. H. Kauffman, Virtual crossing number and the arrow polynomial,
  J. Knot Theory Ramifications, 18(10)(2009), 1335–1357.
- [5] H. A. Dye and L. H. Kauffman, Virtual homotopy, J. Knot Theory Ramifications, 19(7)(2010), 935–960.
- [6] M. Goussarov, M. Polyak and O. Viro, Finite type invariants of classical and virtual knots, Topology, 39(2000), 1045–1068.
- [7] C. Hayashi, A lower bound for the number of Reidemeister moves for unknotting, J. Knot Theory Ramifications, **15(3)**(2006), 313–325.
- [8] M.-J. Jeong, A lower bound for the number of forbidden moves to unknot a long virtual knot, J. Knot Theory Ramifications, **22(6)**(2013), 1350024, 13 pp.
- [9] M.-J. Jeong, *Delta moves and Kauffman polynomials of virtual knots*, J. Knot Theory Ramifications, **23(10)**(2014), 1450053, 17 pp.
- [10] M.-J. Jeong, Reidemeister moves and a polynomial of virtual knot diagrams, J. Knot Theory Ramifications, **24(2)**(2015), 1550010, 16 pp.
- [11] M.-J. Jeong and C.-Y. Park, Similarity indices and the Miyazawa polynomials of virtual links, J. Knot Theory Ramifications, 23(7)(2014), 1460003, 17 pp.
- [12] L. H. Kauffman, Virtual knot theory, Europ. J. Combin., **20**(1999), 663–690.
- [13] L. H. Kauffman, An affine index polynomial invariant of virtual knots, J. Knot Theory Ramifications, **22(4)**(2013), 1340007, 30 pp.
- [14] T. Kanenobu, Forbidden moves unknot a virtual knot, J. Knot Theory Ramifications, 10(1)(2001), 89–96.
- [15] A. Kawauchi, On the Alexander polynomials of knots with Gordian distance one, Topology Appl., 159(2012), 948–958.

- [16] Y. Miyazawa, Magnetic graphs and an invariant for virtual links, J. Knot Theory Ramifications, 15(10)(2006), 1319–1334.
- [17] Y. Miyazawa, A multi-variable polynomial invariant for virtual knots and links, J. Knot Theory Ramifications, 17(11)(2008), 1311–1326.
- [18] Y. Miyazawa, Gordian distance and polynomial invariants, J. Knot Theory Ramifications, 20(6)(2011), 895–907.
- [19] Y. Nakanishi, Unknotting numbers and knot diagrams with the minimum crossings, Math. Sem. Notes Kobe Univ., 11(1983), 257–258.
- [20] S. Nelson, Unknotting virtual knots with Gauss diagram forbidden moves, J. Knot Theory Ramifications, 10(6)(2001), 931–935.
- [21] M. Polyak, Minimal generating sets of Reidemeister moves, Quantum Topol., 1(4)(2010), 399–411.